MTH6115 Cryptography Exercises 3 Solutions

Q1 (a) The output sequence is

Its period is 7.

(b) Every configuration in the same cycle as 1010 will have period 7. The output sequence when the initial configuration is 0001 is

0001011 0001...

This also has period 7 and every configuration in the same cycle as 0001 will have period 7.

The output sequence when the initial configuration is 0000 or 1111 has period 1.

Q2(a)(i) We have $v_0 = u_1, v_1 = u_2, \dots, v_{n-2} = u_{n-1}$ and $v_{n-1} = \sum_{i=0}^{n-1} a_i u_i$. Thus $u_1 = v_0, u_2 = v_1, \dots, u_{n-1} = v_{n-2}$ and

$$u_0 = v_{n-1} + \sum_{i=1}^{n-1} a_i u_i = v_{n-1} + \sum_{i=1}^{n-1} a_i v_{i-1}.$$

(ii) Suppose that $(v_0, v_1, \ldots, v_{n-1})$ is the first configuration which is repeated by the shift register. Part (i) tells us that the configuration $(u_0, u_1, \ldots, u_{n-1})$ which precedes $(v_0, v_1, \ldots, v_{n-1})$ is completely determined by $(v_0, v_1, \ldots, v_{n-1})$. Thus, if $(v_0, v_1, \ldots, v_{n-1})$ were not the initial configuration, then $(u_0, u_1, \ldots, u_{n-1})$ would have been repeated before $(v_0, v_1, \ldots, v_{n-1})$. This would contradict the choice of $(v_0, v_1, \ldots, v_{n-1})$. Hence $(v_0, v_1, \ldots, v_{n-1})$ is the initial configuration. (b) Consider for example the 3-bit shift register described by the polynomial $x^3 + x$, with initial configuration (1, 0, 0). The next two configurations are (0,0,0) and (0,0,0). Thus the first configuration to be repeated is not the initial configuration.

Q3(a) The number of primitive 5-bit shift registers is $\Phi(2^5-1)/5 = \Phi(31)/5 = 6$. (b) Let $p(x) = x^5 + a_4x^4 + a_3x_3 + a_2x^2 + a_1x + a_0$ be irreducible. Since x is not a factor of p(x) we have $a_0 = 1$ and since x + 1 is not a factor we have $a_4 + a_3 + a_2 + a_1 = 1$. This leaves 8 possible values for (a_4, a_3, a_2, a_1) . We can find which of these give irreducible polynomials as follows.

Consider for example $p(x) = x^5 + x^2 + 1$. We know that p(x) has no factor of degree one. So if p(x) factorises, we must have $x^5 + x^2 + 1 = (x^2 + bx + 1)(x^3 + cx^2 + dx + 1)$. Equating powers of x we get b + d = 0, 1 + bd + c = 1, 1 + bc + d = 0, b + c = 0. The first equation gives b = d, and the last equation gives b = c. We can now rewrite the second and third equations as $b^2 + b = 0$ and $b^2 + b = 1$ which clearly have no solution. Hence p(x) is irreducible.

(A similar analysis works for the other irreducible polynomials $x^5 + x^3 + 1$, $x^5 + x^4 + x^3 + x^2 + 1$, $x^5 + x^4 + x^2 + x + 1$, $x^5 + x^4 + x^3 + x + 1$, $x^5 + x^4 + x^3 + x^2 + 1$.)

(b) The output sequence for the shift register described by $x^5 + x^2 + 1$ is

0000100101100100111110001101110101 00001...

The shift register is primitive since it cycles through all 31 non-zero configurations of \mathbb{Z}_2^5 .

Q4 Since the first two letters in the message are 'th', the first 8 bits in the plaintext are 00001001. Since the first 8 bits in the ciphertext are 10111011 So the first 8 bits in the keyword are given by

$0000100101 \oplus 1011101101 = 10110010$

Let the polynomial describing the shift register be $x^4 + a_3x^3 + a_2x^3 + a_1x + a_0$. Then we have the following system of simultaneous equations.

This has the unique solution $a_0 = a_3 = 1$ and $a_1 = a_2 = 0$. So the polynomial describing the shift register is $x^4 + x^3 + 1$. Since the first 8 bits in the keyword are 10110010, the initial configuration of the shift register is (1011). Hence the output sequence of the shift register is

 $101100100011110 \ 1011001000111 \ 10101100100011110 \ 1011001000\ldots$

Since the substitution table is just the addition table for \mathbb{Z}_2 , we can find the plain text by adding the keyword to the ciphertext. Thus the plaintext is

We can now use the table for the International Telegraph Code to decode the plaintext and obtain the message 'threeonetwo'.